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## Data Models and Query Languages Summerterm 2013

### 2. Exercise Sheet: Chase & Datalog

Discussion: 14.05.2013

#### Exercise 1 (Constraints in First-Order Logic)

Consider the following database schema.

```
hasAirport(c_id)  
fly(c_id1,c_id2,dist)  
rail(c_id1,c_id2,dist)
```

Specify the constraints below in First-order Logic and indicate if your specification is a tuple-generating dependency, equality-generating dependency, or none of both. In case of tuple-generating or equality-generating dependencies additionally give their *body* and *head*.

- $\alpha_1$ : If a city has an airport, then there is at least one flight departing from this city.
- $\alpha_2$ : The distance of a rail connection functionally depends on the departure and destination station, i.e. there is only one unique distance for every rail connection.
- $\alpha_3$ : There is at least one flight and one train connection listed in the database.
- $\alpha_4$ : Starting from Frankfurt, all cities with an airport can be reached either by direct flight or by a flight with only one intermediate stop.
- $\alpha_5$ : All pairs of cities with an airport that have a direct train connection also have a direct flight connection.

#### Exercise 2 (Chase Application)

Consider the schema from Exercise 1, the constraint set  $\Sigma := \{\alpha_1, \alpha_2, \alpha_3\}$  with

$$\alpha_1 := \forall c_1, c_2, c_3, d_1, d_2 (\text{rail}(c_1, c_2, d_1), \text{rail}(c_2, c_3, d_2) \rightarrow \exists d_3 \text{rail}(c_1, c_3, d_3))$$
$$\alpha_2 := \forall c_1, c_2, d_1, d_2 (\text{fly}(c_1, c_2, d_1) \wedge \text{fly}(c_2, c_1, d_2) \rightarrow d_1 = d_2)$$
$$\alpha_3 := \forall c_1, c_2, d_1 (\text{fly}(c_1, c_2, d_1) \rightarrow \exists d_2 \text{fly}(c_2, c_1, d_2))$$

and the Conjunctive Query

$$Q: \text{ans}(C_3) \leftarrow \text{rail}(\text{Freiburg}, C_1, D_1), \text{rail}(C_1, C_2, D_2), \text{fly}(C_2, C_3, D_3).$$

- a) Describe the semantics of the constraints and the query informally.
- b) Which constraints from  $\Sigma$  are satisfied by  $body(Q)$ ? Does  $body(Q)$  satisfy  $\Sigma$ ?
- c) Chase query  $Q$  with  $\Sigma$ . Provide all intermediate results (= chase steps). Does it hold that  $body(Q^\Sigma) \models \Sigma$ ?

**Exercise 3 (Datalog, Transitive Closure)**

Given a directed graph  $G$  with edge relation  $E(a, b)$ , which means there is an edge from  $a$  to  $b$  in  $G$ .

- a) Give three different Datalog<sup>+</sup> programs which compute the transitive closure of  $G$ .
- b) Let  $k$  be the length of the longest path in  $G$ . Determine the number of iterations which is needed for each version to compute the transitive closure.
- c) Apply the naive algorithm evaluation of the three programs on the database:

$$E(1, 2), E(2, 3), E(3, 4), E(4, 5)$$

**Exercise 4 (Datalog, Equivalence)**

Consider the following two Datalog<sup>+</sup> programs

$$\begin{aligned} \sqcap_1: \quad & Buys(X, Y) \leftarrow Likes(X, Y) \\ & Buys(X, Y) \leftarrow Knows(X, Z), Buys(Z, Y) \end{aligned}$$

$$\begin{aligned} \sqcap_2: \quad & Buys(X, Y) \leftarrow Likes(X, Y) \\ & Buys(X, Y) \leftarrow Knows(X, Z), Likes(Z, Y) \end{aligned}$$

Proove or falsify that  $\sqcap_1 \equiv \sqcap_2$ .

**Exercise 5 (Datalog, Shortest Paths)**

Given a directed **acyclic** graph  $G$  with edge relation  $E(a, b)$ , which means there is an edge from  $a$  to  $b$  in  $G$ . Give a stratified Datalog program which computes the shortest paths in  $G$  wrt. the number of edges.

**Exercise 6 (Datalog, 3-Coloring of Graphs)**

Let  $G = (V, E)$  be a graph with  $V = \{1, 2, 3, 4, 5\}$  and  $E = \{(1, 2), (1, 4), (1, 5), (2, 3), (2, 4), (3, 4), (4, 5)\}$ .

Consider the following Datalog<sup>+</sup> program

$$\begin{aligned} \sqcap: \quad & Color(N, blue) \leftarrow V(N), not\ Color(N, green), not\ Color(N, red) \\ & Color(N, green) \leftarrow V(N), not\ Color(N, blue), not\ Color(N, red) \\ & Color(N, red) \leftarrow V(N), not\ Color(N, green), not\ Color(N, blue) \\ & NonColoring(N) \leftarrow E(N, M), Color(N, C), Color(M, C) \end{aligned}$$

Give a stable model of  $\sqcap$  where  $NonColoring$  is empty, i.e. demonstrate that the given graph is 3-colorable.